

Fully-adaptive Model for Broadcasting with Universal Lists

Saber Gholami

SYNASC 2022

24th International Symposium on Symbolic and Numeric Algorithms for Scientific Computing

September 12-15, 2022

Linz, Austria.

- 1 Introduction
- 2 Preliminaries and Literature Review
- 3 Fully-adaptive Model
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_{m \times n}$

Tori $T_{m \times n}$

Cube Connected

Cycle CCC_d

Graphs with

$B_{cl}(G) < B_{fa}(G)$

Conclusion and
Future works

Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_{m \times n}$

Tori $T_{m \times n}$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_{fa}(G)$

Conclusion and
Future works

- 1 Introduction
- 2 Preliminaries and Literature Review
- 3 Fully-adaptive Model
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

Introduction



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and
Future works

- Computer networks are becoming more popular each day!
- One problem: Propagate a message
- Information dissemination:
 - ◇ Unicasting,
 - ◇ **Broadcasting**,
 - ◇ Multicasting,
 - ◇ ...



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and
Future works

- Broadcasting is the process of distributing a message from a single node (*originator*) to all other nodes of the network,
- Each *call* is performed during one unit of time,
- Several *calls* could be performed in parallel.

- 1 Introduction
- 2 Preliminaries and Literature Review**
- 3 Fully-adaptive Model
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

- The network: $G = (V, E)$, originator $u \in V$.

- The network: $G = (V, E)$, originator $u \in V$.
- $B_{cl}(u, G)$: minimum time required to finish the broadcasting from u .
- $B_{cl}(G) = \max\{B_{cl}(u, G) | u \in V(G)\}$
 - ◇ For any graph: $B_{cl}(G) \geq \lceil \log n \rceil$

Preliminaries - Classical broadcasting



Introduction

Preliminaries and Literature Review

Fully-adaptive Model

Results on the Fully-adaptive Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and Future works

- The network: $G = (V, E)$, originator $u \in V$.
- $B_{cl}(u, G)$: minimum time required to finish the broadcasting from u .
- $B_{cl}(G) = \max\{B_{cl}(u, G) | u \in V(G)\}$
 - ◊ For any graph: $B_{cl}(G) \geq \lceil \log n \rceil$
- NP-Complete in arbitrary graphs,
- NP-Hard to approximate within a ratio of $(3 - \epsilon)$ for any $\epsilon > 0$.
- Solved optimally for only a few networks,
- A long list of heuristics and approximation algorithms!

- Broadcast scheme: the ordering of the neighbors of each vertex, depending on the originator:
 - ◇ Fix vertex u as the originator,
 - ◇ once vertex v gets informed, it will follow its list I_v^u ,

- Broadcast scheme: the ordering of the neighbors of each vertex, depending on the originator:
 - ◇ Fix vertex u as the originator,
 - ◇ once vertex v gets informed, it will follow its list I_v^u ,
- Drawbacks:
 - ◇ Each vertex has to maintain up to $|V|$ different lists for different originators,
 - ◇ A vertex has to know the originator to perform broadcasting.
 - ◇ Requires a comprehensive knowledge of the network for every vertex.

- Broadcast scheme: the ordering of the neighbors of each vertex, depending on the originator:
 - ◇ Fix vertex u as the originator,
 - ◇ once vertex v gets informed, it will follow its list I_v^u ,
- Drawbacks:
 - ◇ Each vertex has to maintain up to $|V|$ different lists for different originators,
 - ◇ A vertex has to know the originator to perform broadcasting.
 - ◇ Requires a comprehensive knowledge of the network for every vertex.
 - ◇ Inefficient in real-world networks:
 - ◇ Increased message bits,
 - ◇ Need for larger local memory.

- Broadcasting with universal lists:
 - ◇ Each vertex v has a single list l_v to follow, regardless of the originator.

- [1] Slater, P.J., Cockayne, E.J. and Hedetniemi, S.T., 1981. Information dissemination in trees. *SIAM Journal on Computing*, 10(4), pp.692-701..
- [2] Diks, K. and Pelc, A., 1996. Broadcasting with universal lists. *Networks*, 27(3), pp.183-196.

Preliminaries - Universal lists broadcasting



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_s(G)$

Conclusion and
Future works

- Broadcasting with universal lists:
 - ◇ Each vertex v has a single list l_v to follow, regardless of the originator.
- Two sub-models:
 - ◇ Non-adaptive $B_{na}(G)$: send to all vertices on the list,
 - ◇ Adaptive $B_a(G)$: skip the vertices from which the message is received.

[1] Slater, P.J., Cockayne, E.J. and Hedetniemi, S.T., 1981. Information dissemination in trees. *SIAM Journal on Computing*, 10(4), pp.692-701..

[2] Diks, K. and Pelc, A., 1996. Broadcasting with universal lists. *Networks*, 27(3), pp.183-196.

Preliminaries - Universal lists broadcasting



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees
Grids $G_m \times n$
Tori $T_m \times n$
Cube Connected
Cycle CCC_d
Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

- Broadcasting with universal lists:
 - ◇ Each vertex v has a single list l_v to follow, regardless of the originator.
- Two sub-models:
 - ◇ Non-adaptive $B_{na}(G)$: send to all vertices on the list,
 - ◇ Adaptive $B_a(G)$: skip the vertices from which the message is received.
- Introduced indirectly by Slater [1]; for any Tree, $B_{cl}(T) = B_a(T)$.
- Diks and Pelc [2] distinguished between adaptive and non-adaptive models,
 - ◇ Also proposed several broadcast schemes for different graphs
- Long list of research ...

[1] Slater, P.J., Cockayne, E.J. and Hedetniemi, S.T., 1981. Information dissemination in trees. *SIAM Journal on Computing*, 10(4), pp.692-701..

[2] Diks, K. and Pelc, A., 1996. Broadcasting with universal lists. *Networks*, 27(3), pp.183-196.

Introduction

Preliminaries
and Literature
Review

**Fully-adaptive
Model**

Results on the
Fully-adaptive
Model

Trees

Grids $G_{m \times n}$

Tori $T_{m \times n}$

Cube Connected

Cycle CCC_d

Graphs with

$B_{cl}(G) < B_{fa}(G)$

Conclusion and
Future works

- 1 Introduction
- 2 Preliminaries and Literature Review
- 3 Fully-adaptive Model**
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

Fully-adaptive Model



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_s(G)$

Conclusion and
Future works

- Another sub-model for universal lists,
- A universal list l_u is maintained at each vertex u ,
- Once informed, follow the list and **skip all informed vertices!**
 - ◊ Similarly to the classical model: No unnecessary calls!

Fully-adaptive Model



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

Cube Connected
Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

- Another sub-model for universal lists,
- A universal list l_u is maintained at each vertex u ,
- Once informed, follow the list and **skip all informed vertices!**
 - ◊ Similarly to the classical model: No unnecessary calls!
- **Theorem 1.** $B_{cl}(G) \leq B_{fa}(G) \leq B_a(G) \leq B_{na}(G)$, for any graph G .

Model	Symbol	No. of unnecessary calls	Space Complexity	Speed
Non-adaptive	$B_{na}(G)$	Many	Very Low	Very Slow
Adaptive	$B_a(G)$	Few	Low	Slow
Fully Adaptive	$B_{fa}(G)$	0	Moderate	Moderate
Classical	$B_{cl}(G)$	0	Very High	Very Fast

- A broadcast scheme: Matrix $\sigma_{n \times \Delta}$,
 - ◊ Row i of σ corresponds to an ordering for vertex v_i .
- Set of all possible schemes: Σ .

Fully-adaptive Model - Definitions



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_s(G)$

Conclusion and
Future works

- A broadcast scheme: Matrix $\sigma_{n \times \Delta}$,
 - ◇ Row i of σ corresponds to an ordering for vertex v_i .
- Set of all possible schemes: Σ .
- Let $M \in \{na, a, fa\}$ be a model:
 - ◇ $B_M^\sigma(v, G)$: the time steps needed to inform all the vertices in G from v while following σ under M ,
 - ◇ $B_M^\sigma(G) = \max_{v \in V} \{B_M^\sigma(v, G)\}$,
 - ◇ $B_M(G) = \min_{\sigma \in \Sigma} \{B_M^\sigma(G)\}$.

Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

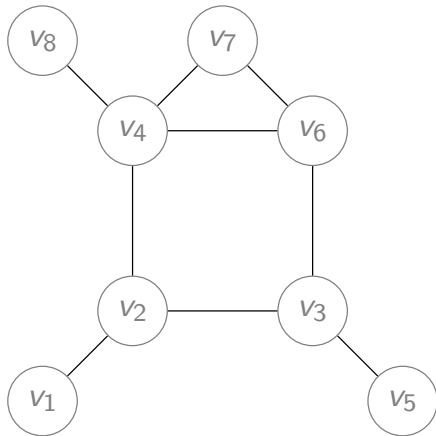
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

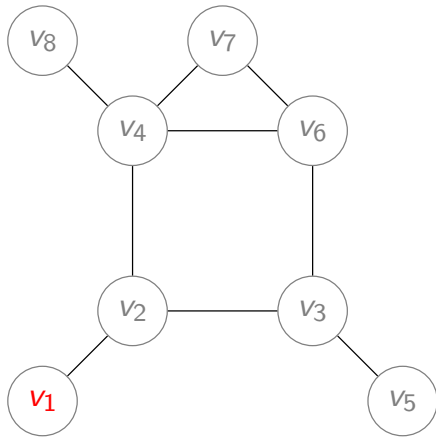
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

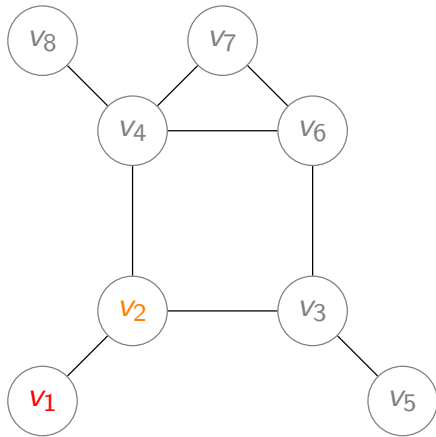
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

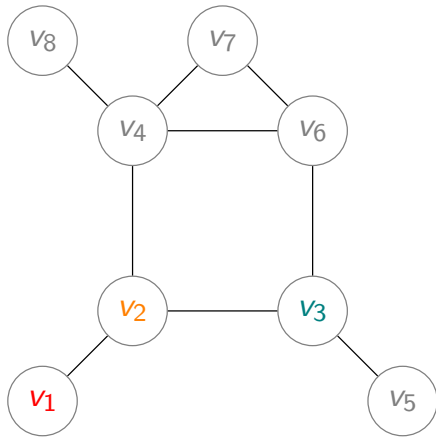
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

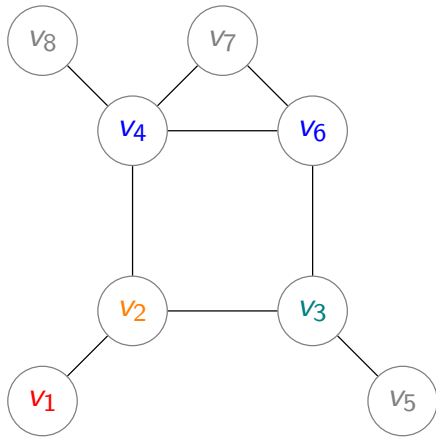
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

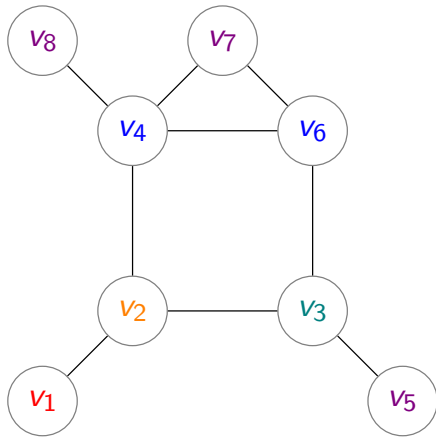
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



Fully-adaptive model - Example



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

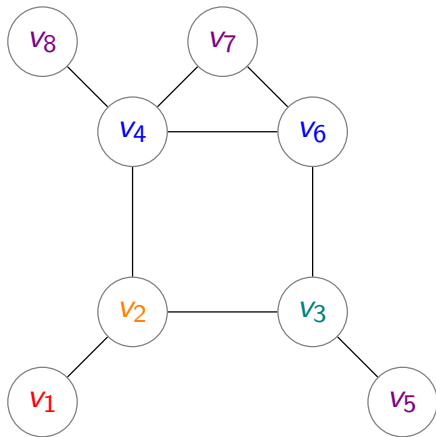
Cube Connected

Cycle CCC_d

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

Sender	Ordering of receivers			
v_1	v_2	Null	Null	Null
v_2	v_3	v_4	v_1	Null
v_3	v_2	v_6	v_5	Null
v_4	v_2	v_6	v_8	v_7
v_5	v_3	Null	Null	Null
v_6	v_3	v_7	v_4	Null
v_7	v_6	v_4	Null	Null
v_8	v_4	Null	Null	Null



- $B_{fa}^\sigma(v_1, G) = 4$, while $B_a^\sigma(v_1, G) = 5$ and $B_{na}^\sigma(v_1, G) = 6$.

- Assumptions:
 - ◇ None-faulty network with established links,
 - ◇ Unique and heavy message,
 - ◇ The message: header + payload,

Fully-adaptive Model - AAA



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and
Future works

- Assumptions:
 - ◇ None-faulty network with established links,
 - ◇ Unique and heavy message,
 - ◇ The message: header + payload,
- Architecture:
 - ◇ How to know the state of each neighbour?
 - ◇ Push model,
 - ◇ Pull model,

Fully-adaptive Model - AAA



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and
Future works

- **Assumptions:**
 - ◇ None-faulty network with established links,
 - ◇ Unique and heavy message,
 - ◇ The message: header + payload,
- **Architecture:**
 - ◇ How to know the state of each neighbour?
 - ◇ Push model,
 - ◇ Pull model,
- **Applications:**
 - ◇ Update procedure in SDNs:
 - ◇ Changing routing policies, adjusting links' weights, etc.
 - ◇ The data plane only forwards packets,
 - ◇ Routing and load balancing decisions are made in a centralized controller,
 - ◇ The network manager must optimize the forwarding tables (broadcast schemes).

- 1 Introduction
- 2 Preliminaries and Literature Review
- 3 Fully-adaptive Model
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_{m \times n}$

Tori $T_{m \times n}$

Cube Connected

Cycle CCC_d

Graphs with

$B_{cl}(G) < B_{fa}(G)$

Conclusion and
Future works

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

- Trees T :

- ◇ **Theorem 2.** $B_{cl}(T) = B_{fa}(T) = B_a(T)$.

- Grids $G_{m \times n}$:
 - ◇ **Corollary 1.**
 $B_{fa}(G_{m \times n}) = m + n - 2.$

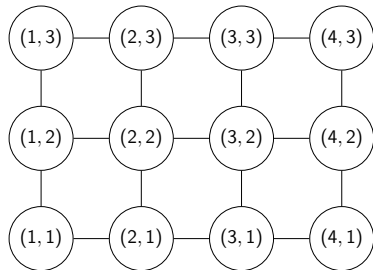


Figure: Grid with $m = 4, n = 3$

Results on the Fully-adaptive Model - cont.



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times n$

Tori $T_m \times n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_b(G)$

Conclusion and
Future works

- Tori $T_{m \times n}$:
 - ◇ **Theorem 3.**
 - ◇ $B_{fa}(T_{m \times n}) = \lfloor \frac{n}{2} \rfloor + \lfloor \frac{m}{2} \rfloor$,
if m and n are even,
 - ◇ $B_{fa}(T_{m \times n}) = \lfloor \frac{n}{2} \rfloor + \lfloor \frac{m}{2} \rfloor + 1$,
if only one of m and n is even,
 - ◇ $\lfloor \frac{n}{2} \rfloor + \lfloor \frac{m}{2} \rfloor + 1 \leq B_{fa}(T_{m \times n}) \leq$
 $\lfloor \frac{n}{2} \rfloor + \lfloor \frac{m}{2} \rfloor + 2$,
if m and n are odd.

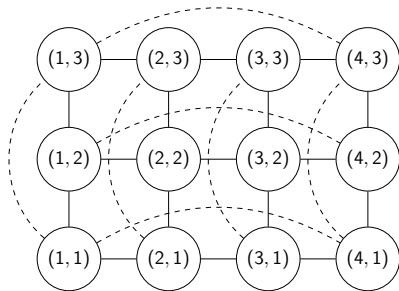


Figure: Torus with $m = 4, n = 3$

Results on the Fully-adaptive Model - cont.



Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

**Cube Connected
Cycle CCC_d**

Graphs with
 $B_d(G) < B_b(G)$

Conclusion and
Future works

- Cube Connected Cycle CCC_d :

- ◇ **Theorem 4.**

- ◇ $B_{fa}(CCC_d) = \lceil \frac{5d}{2} \rceil - 1$

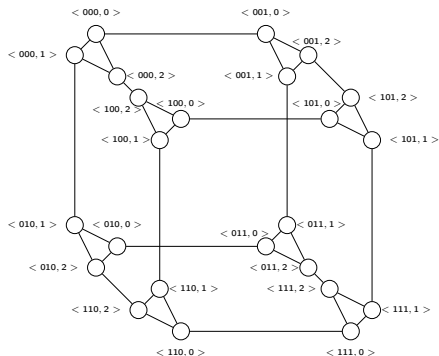


Figure: Cube-Connected Cycle with $d = 3$

Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_m \times_n$

Tori $T_m \times_n$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_{fa}(G)$

Conclusion and
Future works

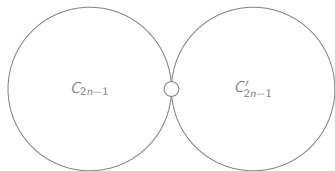
- Is $B_{cl}(G) = B_{fa}(G)$ always?

Results on the Fully-adaptive Model - cont.

- Is $B_{cl}(G) = B_{fa}(G)$ always?

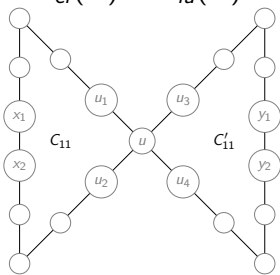
- ◊ No!

- ◊ **Proposition 1.** *There exists graph G with $B_{cl}(G) < B_{fa}(G)$:*



◊

a)



b)

Introduction

Preliminaries
and Literature
Review

Fully-adaptive
Model

Results on the
Fully-adaptive
Model

Trees

Grids $G_{m \times n}$

Tori $T_{m \times n}$

Cube Connected

Cycle CCC_d

Graphs with

$B_d(G) < B_{fa}(G)$

Conclusion and
Future works

- 1 Introduction
- 2 Preliminaries and Literature Review
- 3 Fully-adaptive Model
- 4 Results on the Fully-adaptive Model
 - Trees
 - Grids $G_{m \times n}$
 - Tori $T_{m \times n}$
 - Cube Connected Cycle CCC_d
 - Graphs with $B_{cl}(G) < B_{fa}(G)$
- 5 Conclusion and Future works

- We proposed the Fully-adaptive model for broadcasting with universal lists
- Benefits:
 - ◊ Uses less memory compared to classical broadcasting,
 - ◊ Faster than the adaptive and non-adaptive models,
- We designed optimal broadcast schemes for:
 - ◊ Trees,
 - ◊ Grids,
 - ◊ Tori,
 - ◊ Cube-connected Cycle,
- We designed graphs with $B_{cl}(G) < B_{fa}(G)$.

- Finding the broadcast time of other families of graphs under the fully-adaptive model, such as:
 - ◇ Complete graph,
 - ◇ Shuffle-exchange,
 - ◇ DeBruijn,
 - ◇ Complete bipartite graph,
 - ◇ etc.
- Designing an efficient algorithm to find $B_{fa}(G)$ for general graphs, or prove that it is NP-Hard.

